

# BLUE WATERS

SUSTAINED PETASCALE COMPUTING

## Performance Modeling for Systematic Performance Tuning

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GREAT LAKES CONSORTIUM  
FOR PETASCALE COMPUTATION

## Special Announcement!

- Blue Waters is now officially back!



- ... but back to the talk (examples are still POWER7)

Details: <http://www.ncsa.illinois.edu/BlueWaters/system.html>

## The Perspective of a Computing Center

- Performance = “completed science per cost and time”
- Optimizing this metric can be manifold:
  - Application optimization (support application teams)
  - Architecture optimization (select best hardware)
  - Optimize Middleware (scheduler, libraries etc.)
  - Optimize Policies (scheduling, charging etc.)
  - ... and many more

## Performance Modeling – State of the Practice

- Delivers the “science per cost/time” metric
  - Can be used to drive optimizations!
- Who does performance modeling?
  - Mostly computer scientists, in-house teams
- BUT: most development is done by application developers and/or domain scientists
  - They should develop performance models during software development
    - See performance modeling panel @3:30 in TCC 101

## (Ideal) State of the Practice @NCSA

- Propose to use simple performance modeling to characterize the behavior of applications
  - Enables rough optimization (cf. “80/20 rule”)
- We provide a set of simple modeling guidelines
  - Semi-analytic performance modeling
  - Small number of parameters, use other techniques where necessary

Benchmark    ---- Full Simulation    ---- Model Simulation    ---- Model

Number of Parameters

Model Error

## Overview of Performance Modeling

- Analytic modeling:
  - Determine application requirements and system speeds to compute time (e.g., bandwidth)
- Empirical modeling (e.g. [1,2]):
  - “Black-box” approach: machine learning, neural networks, statistical learning ...
- Semi-empirical modeling:
  - “White box” approach: find asymptotically tight analytic models, parameterize empirically (curve fitting)

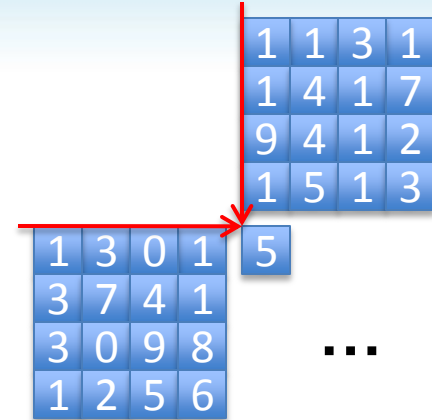


[1]: Barnes, Rountree, Lowenthal, Reeves, Supinski, Schulz: A regression-based approach to scalability prediction  
[2]: McKee, Singh, Supinski, Schulz: Constructing Application Performance Models Using Neural Networks

## A Quick Example - MM

- Matrix multiplication ( $N^3$  algorithm)

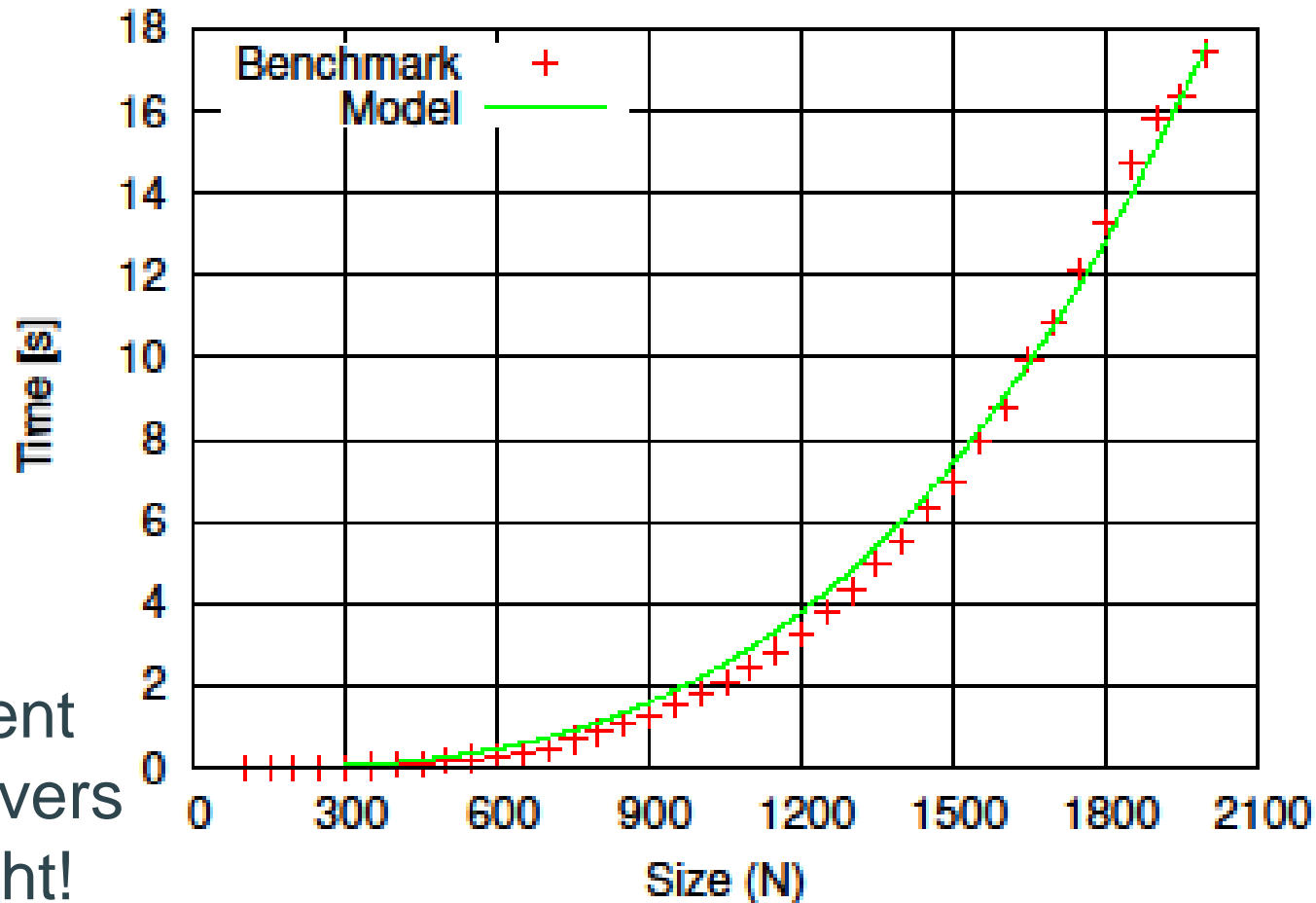
```
for(int i=0; i<N; ++i)
  for(int j=0; j<N; ++j)
    for(int k=0; k<N; ++k)
      C[i+j*N] += A[i+k*N] * B[k+j*N];
```



- Trivial (non-blocked) algorithm
- Analytic Model:
  - $N^3$  FP add/mult,  $4N^3$  FP load/store, +int ops
  - How can we get to an execution time? → **very hard!**

## Semi-Empiric Model for MM

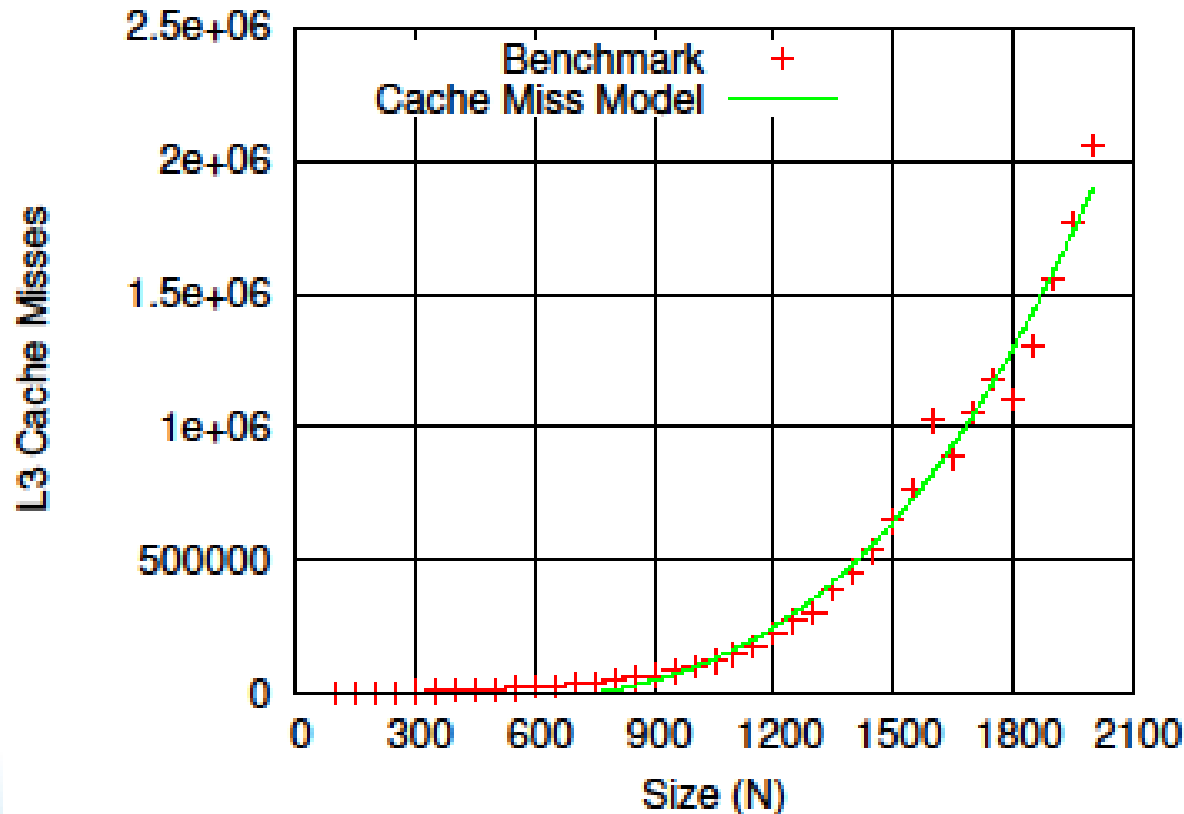
- $T(N) = tN^3$
- POWER7
  - $t=2.2\text{ns}$
  - 0.8% err
- Is that all?
  - Requirement Model delivers more insight!





# Requirements Model for MM

- Required floating point operations:  $2N^3$  (verified)
- Cache misses?
  - Semi-analytic!
  - $C(N) = aN^3 - bN^2$
- POWER7
  - $a=3.8e-4$
  - $a=2.7e-1$



## Our Ubiquitous Modeling Philosophy

- Modeling during each phase of SW development:
  - Analysis – pick right method (asymptotic models)
  - Design – pick right algorithms (asymptotic models)
  - Implementation – show good usage of machine, e.g., blocking in MM (semi-empirical models)
  - Testing – fulfilling model expectations as correctness criterion (compare tests with models)
  - Maintenance – monitor performance on different architectures (compare times with models)

## More uses of Models

- Performance Optimization
  - Identify bottlenecks and problems during porting
- System Design
  - Co-design based on application requirements
- System Deployment and Testing
  - Know what to expect, find performance issues quickly
- During System Operation
  - Detect silent (and slow) performance degradation



## Six-Steps to a Model

- Our very high-level strategy consists of the following six steps:

- 1) Identify input parameters that influence runtime
- 2) Identify application kernels
- 3) Determine communication pattern
- 4) Determine communication/computation overlap

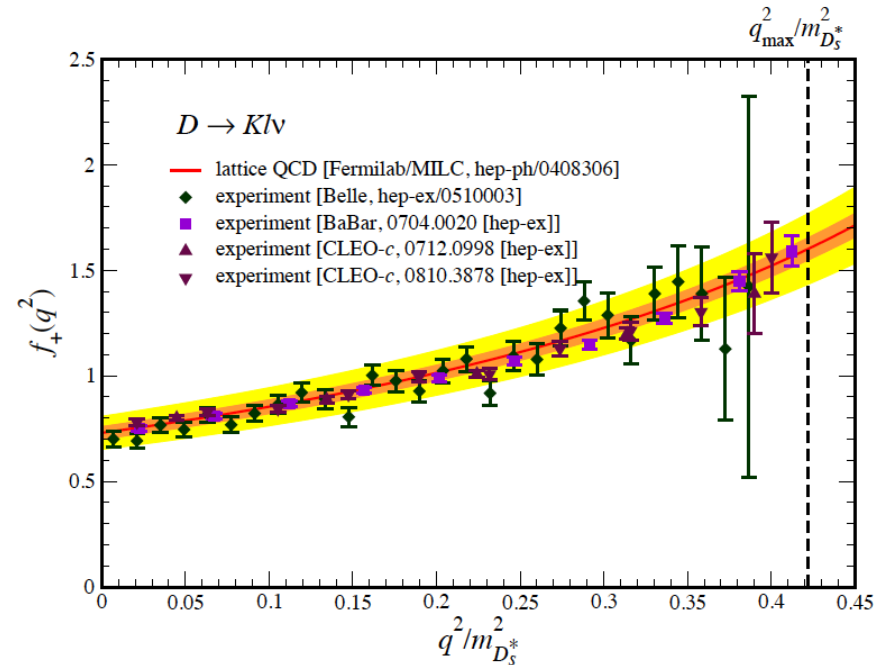
Analytic

- 5) Determine sequential baseline
- 6) Determine communication parameters

Empiric

# All Steps By Example – MILC

- MIMD Lattice Computation
  - Gains deeper insights in fundamental laws of physics
  - Determine the predictions of lattice field theories (QCD & Beyond Standard Model)
  - Major NSF application
- Challenge:
  - High accuracy (computationally intensive) required for comparison with results from experimental programs in high energy & nuclear physics



Bernard, Gottlieb et al.: Studying Quarks and Gluons On Mimd Parallel Computers

# Step 1: Critical Parameters

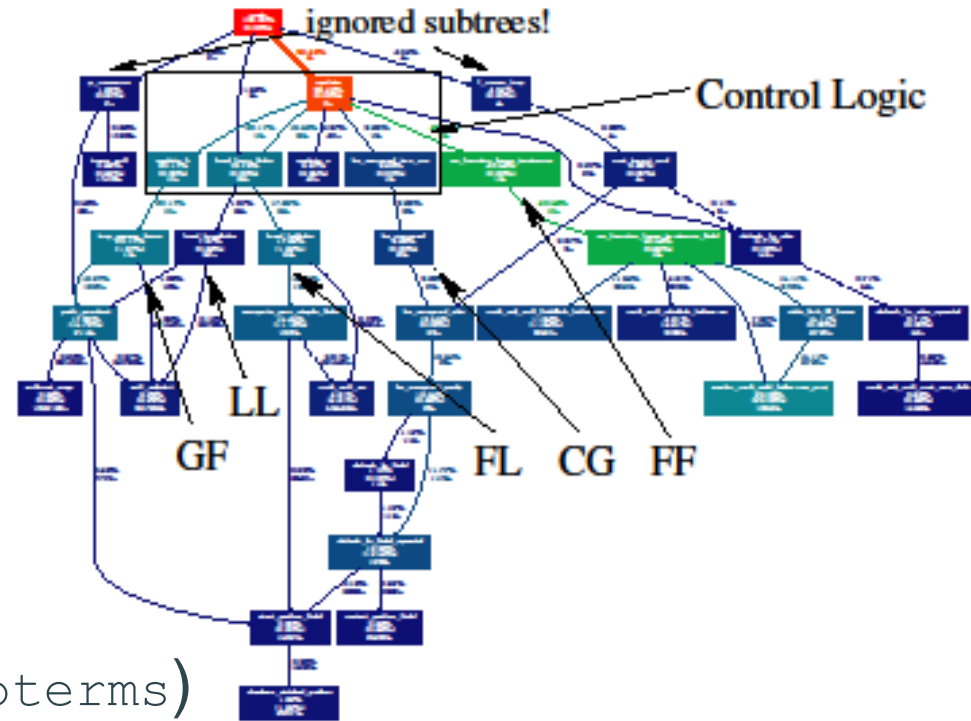
- Best way: ask a domain expert!
  - Or: look through the code/input file format
- For MILC (thanks to S. Gottlieb):



Name	Description
P	number of PEs (intrinsic parameter)
nx, ny, nz, nt	size in x, y, z, t dimension
warms, trajecs	warmup rounds and trajectories (outer loop)
traj_between_meas	measurement “frequency”
steps_per_trajectory	number of “steps” in each trajectory
beta, mass1, ...	physics parameters that influence CG iterations
max_cg_iterations	limits the conjugate gradient iterations

## Step 2: Find Kernels

- E.g., investigate call-tree or source-code
- Control logic
  - update
- MILC's kernels:
  - LL (load\_longlinks)
  - FL (load\_fatlinks)
  - CG (ks\_congrad)
  - GF (imp\_gauge\_force)
  - FF (eo\_fermion\_force\_twoterms)



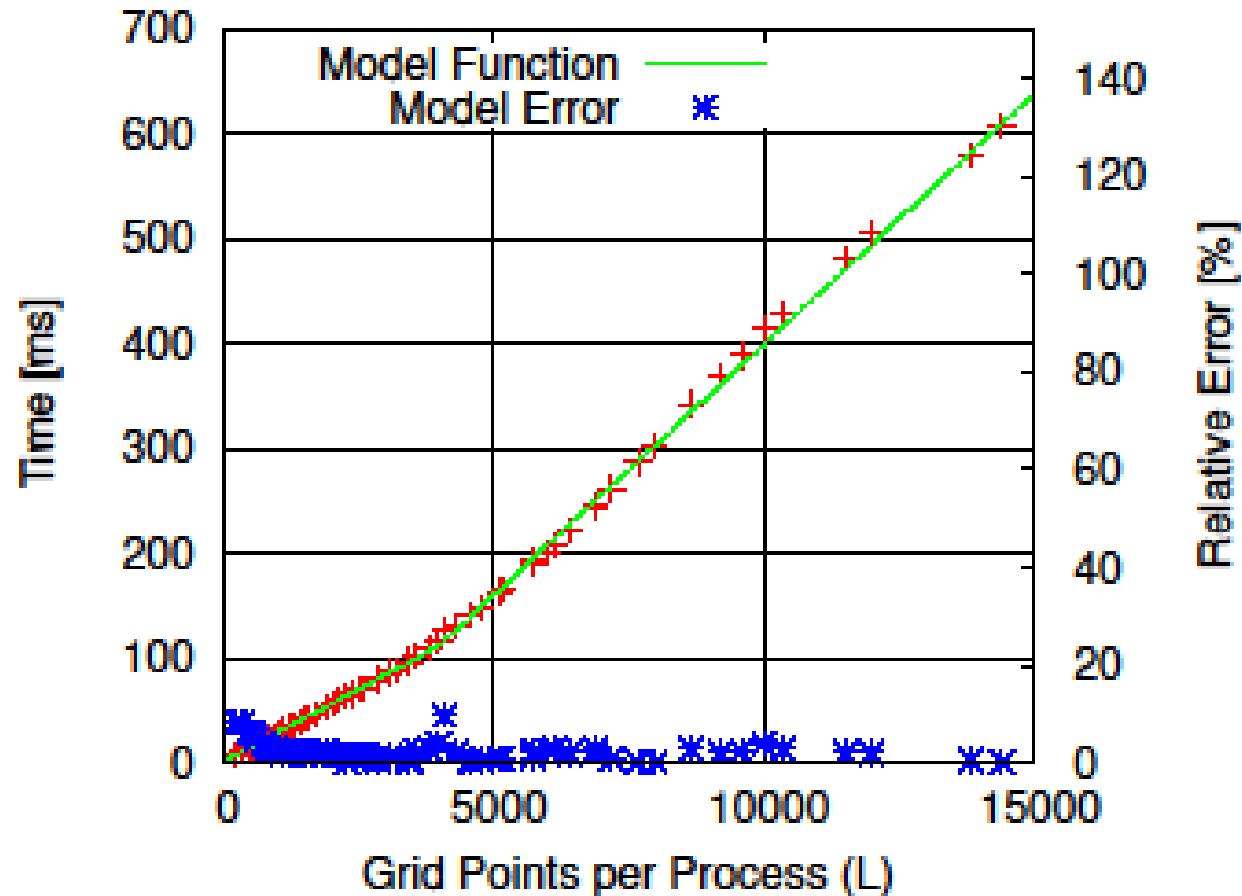
## Step 4: Sequential Performance

- MILC “only” loops over the lattice  $\rightarrow \Theta(V)$
- $T(V) = tV$ 
  - Wait, it’s not that simple with caches ☹
  - Small  $V$  fit in cache!
- $T(V) = t_1 * \min(s, V) + t_2 * \max(0, V-s)$ 
  - Cache holds  $s$  data elements
  - Three parameters for each kernel



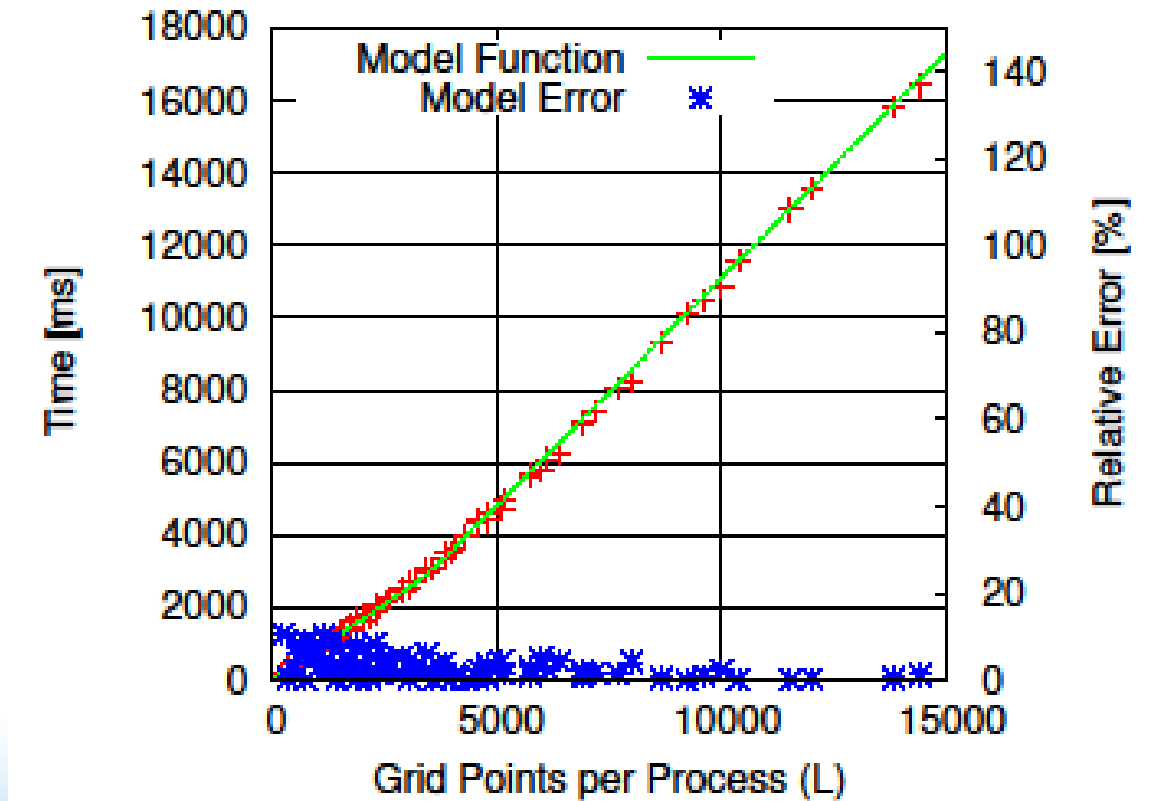
## An Example Kernel: GF (Gauge Force)

- On POWER7:
  - $t_1=62.4 \mu\text{s}$
  - $t_2=92 \mu\text{s}$
  - $s=4.000$
- Errors
  - Max <10%
  - Cum <3%



# Complete Serial Performance Model

$$T_{serial}(V) = (\text{trajec} + \text{warms}) \cdot \text{steps} \cdot [T(FF, V) + T(GF, V) + 3(T(LL, V) + T(FL, V))] + \left[ \frac{\text{trajec}}{\text{meas}} \right] [T(LL, V) + T(FL, V)] + \text{niters} \cdot T(CG, V)$$



## Step 3: Communication Pattern

- 4d domain is cut in all dimensions (cubic)
  - 4d nearest-neighbor communication (8 neighbors)
- Allreduce to check CG convergence
  - One per iteration on full process set
- We counted messages and sizes
  - Separate for each kernel
  - See paper for sizes and full model equation!

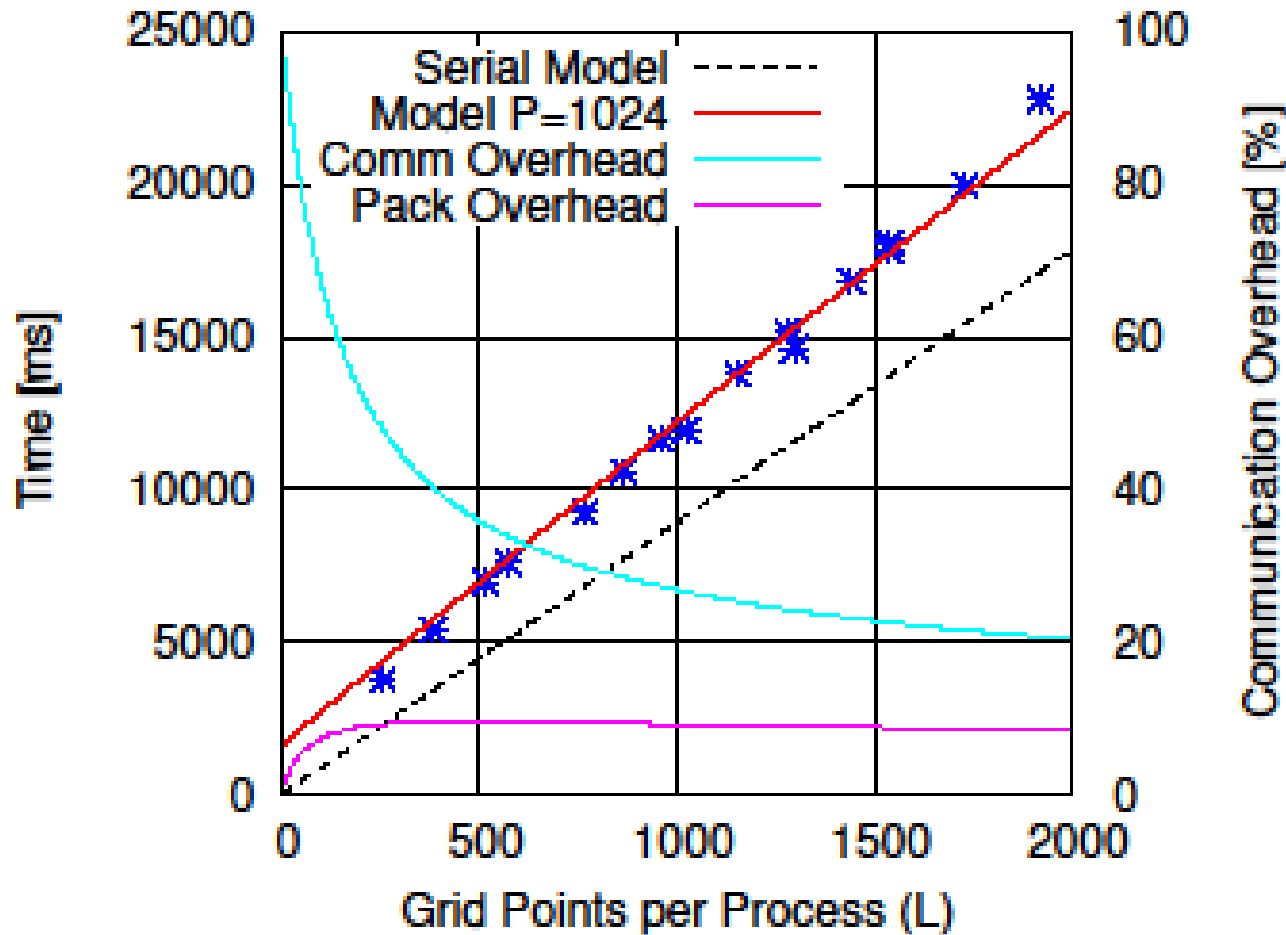
kernel	$\#Messages$
FF	$(trajec s + warms) \cdot steps \cdot 1616$
GF	$(trajec s + warms) \cdot steps \cdot 828$
LL	$(3 \cdot steps \cdot (trajec s + warms) + \lfloor \frac{trajec s}{meas} \rfloor) \cdot 8$
FL	$(3 \cdot steps \cdot (trajec s + warms) + \lfloor \frac{trajec s}{meas} \rfloor) \cdot 288$

## Step 6: Communication Parameters

- Two options:
  - Semi-empiric – fit measurements to get effective latency and bandwidth
    - Enables to check if they match expectations
  - Analytic – derive parameters separately (e.g., documentation or separate benchmark)
    - Often problematic if they do not match expectations
- Our model was analytic
  - Uses LogGP parameters (measured by Netgauge [1])

[1] Hoefler et al.: Low-Overhead LogGP Parameter Assessment for Modern Interconnection Networks

# The Fully-Parameterized Parallel Model



## Conclusions and Future Work

- Models in use for predictions and optimizations
  - First successes: ~10-20% improved performance [1]
- Simple strategy enables application team models
  - Better chance to be maintained than external models
  - Critical for performance-centric software development
- We need (and work on):
  - More examples for irregular/dynamic codes
  - Better tool support for modeling



[1] Hoefler, Gottlieb.: Parallel Zero-Copy Algorithms for Fast Fourier Transform and Conjugate Gradient using MPI Datatypes